

Proofs for: Construction Principles for Well-behaved Scalable Systems

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Theorem 1 (intersection theorem). *Let \mathcal{I} be a parameter structure, $\mathcal{B}_{\mathcal{I}}$ an isomorphism structure for \mathcal{I} , and $T \neq \emptyset$.*

- i) *Let $(\mathcal{L}_I^t)_{I \in \mathcal{I}}$ for each $t \in T$ be a monotonic parameterised system, then $(\bigcap_{t \in T} \mathcal{L}_I^t)_{I \in \mathcal{I}}$ is a monotonic parameterised system.*
- ii) *Let $(\mathcal{L}_I^t)_{I \in \mathcal{I}}$ for each $t \in T$ be a scalable system with respect to $\mathcal{B}_{\mathcal{I}}$, then $(\bigcap_{t \in T} \mathcal{L}_I^t)_{I \in \mathcal{I}}$ is a scalable system with respect to $\mathcal{B}_{\mathcal{I}}$.*
- iii) *Let $(\mathcal{L}_I^t)_{I \in \mathcal{I}}$ for each $t \in T$ be a self-similar monotonic parameterised system, then $(\bigcap_{t \in T} \mathcal{L}_I^t)_{I \in \mathcal{I}}$ is a self-similar monotonic parameterised system.*

Proof of Theorem 1 (i)–(iii):

Proof of (i): Let $(\mathcal{L}_I^t)_{I \in \mathcal{I}}$ a monotonic parameterised system for each $t \in T$, then $\mathcal{L}_{I'}^t \subset \mathcal{L}_I^t$ for $t \in T$, $I, I' \in \mathcal{I}$, and $I' \subset I$. This implies

$$\bigcap_{t \in T} \mathcal{L}_{I'}^t \subset \bigcap_{t \in T} \mathcal{L}_I^t.$$

So, $(\bigcap_{t \in T} \mathcal{L}_I^t)_{I \in \mathcal{I}}$ is a monotonic parameterised system.

Proof of (ii): Let $(\mathcal{L}_I^t)_{I \in \mathcal{I}}$ an scalable system with respect to $(\mathcal{B}(I, K))_{(I, K) \in \mathcal{I} \times \mathcal{I}}$ for each $t \in T$, then $\iota_K^I(\mathcal{L}_I^t) = \mathcal{L}_K^t$ for $t \in T$, $I, K \in \mathcal{I}$, and $\iota \in \mathcal{B}(I, K)$.

Because all ι_K^I are isomorphisms,

$$\iota_K^I(\bigcap_{t \in T} \mathcal{L}_I^t) = \bigcap_{t \in T} \iota_K^I(\mathcal{L}_I^t) = \bigcap_{t \in T} \mathcal{L}_K^t.$$

Therefore, $(\bigcap_{t \in T} \mathcal{L}_I^t)_{I \in \mathcal{I}}$ is a scalable system

with respect to $(\mathcal{B}(I, K))_{(I, K) \in \mathcal{I} \times \mathcal{I}}$.

Proof of (iii): Let $(\mathcal{L}_I^t)_{I \in \mathcal{I}}$ a self-similar monotonic parameterised system for each $t \in T$. For $I, I' \in \mathcal{I}$ with $I' \subset I$ holds

$$\Pi_{I'}^I(\bigcap_{t \in T} \mathcal{L}_I^t) \subset \bigcap_{t \in T} \Pi_{I'}^I(\mathcal{L}_I^t) = \bigcap_{t \in T} \mathcal{L}_{I'}^t \subset \bigcap_{t \in T} \mathcal{L}_I^t. \quad (1)$$

Because $\bigcap_{t \in T} \mathcal{L}_{I'}^t \subset \Sigma_{I'}^*$ holds

$$\Pi_{I'}^I(\bigcap_{t \in T} \mathcal{L}_I^t) = \bigcap_{t \in T} \mathcal{L}_{I'}^t.$$

Together with the second inclusion from (1) it follows

$$\bigcap_{t \in T} \mathcal{L}_{I'}^t \subset \Pi_{I'}^I(\bigcap_{t \in T} \mathcal{L}_I^t).$$

Because of the first part of (1) now holds

$$\Pi_{I'}^I(\bigcap_{t \in T} \mathcal{L}_I^t) = \bigcap_{t \in T} \mathcal{L}_{I'}^t.$$

Therefore,

$$(\bigcap_{t \in T} \mathcal{L}_I^t)_{I \in \mathcal{I}}$$

is a self-similar monotonic parameterised system with respect to \mathcal{I} . ■

Theorem 2 (simplest well-behaved scalable systems). *$(\dot{\mathcal{L}}(L)_I)_{I \in \mathcal{I}}$ is a well-behaved scalable system with respect to each isomorphism structure for \mathcal{I} based on N and*

$$\dot{\mathcal{L}}(L)_I = \bigcap_{i \in N} (\tau_i^I)^{-1}(L) \text{ for each } I \in \mathcal{I}.$$

The proof of Theorem 2 will be given in context of influence structures because it consists of special cases of more general results on influence structures (see 32).

Further requirements, which assure that $(\mathcal{L}(L, \mathcal{E}_{\mathcal{I}}, V)_I)_{I \in \mathcal{I}}$ are well-behaved scalable systems, will be given with respect to $\mathcal{E}_{\mathcal{I}}$, $\mathcal{B}_{\mathcal{I}}$, L and V . This will be prepared by some lemmata.

Lemma 1. *Let $\mathcal{E}_{\mathcal{I}} := (E(t, I))_{(t, I) \in T \times \mathcal{I}}$ be an influence structure for \mathcal{I} indexed by T , and let $V \subset \Sigma^*$. If*

$$E(t, I') = E(t, I) \cap I' \quad (2)$$

for each $t \in T$ and $I, I' \in \mathcal{I}$ $I' \subset I$, then

$$((\tau_{E(t, I)})^{-1}(V))_{I \in \mathcal{I}}$$

is a monotonic parameterised system for each $t \in T$, and by the intersection theorem

$$(\bigcap_{t \in T} (\tau_{E(t, I)})^{-1}(V))_{I \in \mathcal{I}}$$

is a monotonic parameterised system.

Proof: Let $I \in \mathcal{I}$ and $t \in T$. From the definitions of influence homomorphisms and influence structures it follows

$$\tau_{E(t,I)}^I(a_i) = \begin{cases} a & | & a_i \in \Sigma_{E(t,I)} \\ \varepsilon & | & a_i \in \Sigma_I \setminus \Sigma_{E(t,I)} \end{cases} . \quad (3)$$

For $I' \subset I$, $I' \in \mathcal{I}$ and $a_i \in \Sigma_{I'}$ then because of (2)

$$\begin{aligned} \tau_{E(t,I)}^I(a_i) &= \begin{cases} a & | & a_i \in \Sigma_{E(t,I)} \cap \Sigma_{I'} \\ \varepsilon & | & a_i \in \Sigma_{I'} \cap \Sigma_I \setminus \Sigma_{E(t,I)} \end{cases} \\ &= \begin{cases} a & | & a_i \in \Sigma_{E(t,I')} \\ \varepsilon & | & a_i \in \Sigma_{I'} \setminus (\Sigma_{E(t,I)} \cap \Sigma_{I'}) \end{cases} \\ &= \begin{cases} a & | & a_i \in \Sigma_{E(t,I')} \\ \varepsilon & | & a_i \in \Sigma_{I'} \setminus \Sigma_{E(t,I')} \end{cases} = \tau_{E(t,I')}^{I'}(a_i), \end{aligned}$$

and therefore

$$(\tau_{E(t,I')}^{I'})^{-1}(V) \subset (\tau_{E(t,I)}^I)^{-1}(V) \text{ for } V \subset \Sigma^*. \quad (4)$$

So,

$$((\tau_{E(t,I)}^I)^{-1}(V))_{I \in \mathcal{I}} \quad (5)$$

is a monotonic parameterised system for each $t \in T$. ■

Example 1. Let \mathcal{I} be a parameter structure based on N . For $I \in \mathcal{I}$ and $i \in N$ let:

$$\dot{E}(i, I) := \begin{cases} \{i\} & | & i \in I \\ \emptyset & | & i \in N \setminus I \end{cases} .$$

By the definition of parameter structure $N \neq \emptyset$. So

$$\dot{\mathcal{E}}_{\mathcal{I}} := (\dot{E}(i, I))_{(i, I) \in N \times \mathcal{I}}$$

defines an influence structure for \mathcal{I} indexed by N . $\dot{\mathcal{E}}_{\mathcal{I}}$ satisfies (2) and by $\tau_i^I = \tau_{\{i\}}^I$ $\tau_i^I = \tau_{\dot{E}(i, I)}^I$ for $i \in N$ and $I \in \mathcal{I}$.

Now by Lemma 1 for $V \subset \Sigma^*$

$$((\tau_i^I)^{-1}(V))_{I \in \mathcal{I}} \text{ is a monotonic parameterised system} \quad (6)$$

for each $i \in N$.

For this special influence structure $\dot{\mathcal{E}}_{\mathcal{I}}$ a stronger result can be obtained.

Lemma 2. Let \mathcal{I} be a parameter structure based on N and $\varepsilon \in L \subset \Sigma^*$. Then

$$((\tau_i^I)^{-1}(L))_{I \in \mathcal{I}}$$

is a self-similar monotonic parameterised system for each $i \in N$, and by the intersection theorem

$$\left(\bigcap_{i \in N} (\tau_i^I)^{-1}(L) \right)_{I \in \mathcal{I}}$$

is a self-similar monotonic parameterised system.

Proof: On account of (6)

$$\Pi_{I'}^I((\tau_i^I)^{-1}(L)) = (\tau_i^{I'})^{-1}(L)$$

has to be shown for $I, I' \in \mathcal{I}$, $I' \subset I$, and $i \in N$.

(6) implies $(\tau_i^{I'})^{-1}(L) \subset (\tau_i^I)^{-1}(L)$ and therefore,

$$(\tau_i^{I'})^{-1}(L) = \Pi_{I'}^I((\tau_i^I)^{-1}(L)) \subset \Pi_{I'}^I((\tau_i^I)^{-1}(L)). \quad (7)$$

It remains to show $\Pi_{I'}^I((\tau_i^I)^{-1}(L)) \subset (\tau_i^{I'})^{-1}(L)$.

Case 1. $i \notin I'$

Because of $\varepsilon \in L$ and $\tau_i^{I'}(w) = \varepsilon$ for $i \notin I'$ and $w \in \Sigma_{I'}^*$, it holds $(\tau_i^{I'})^{-1}(L) = \Sigma_{I'}^*$ and so

$$\Pi_{I'}^I((\tau_i^I)^{-1}(L)) \subset (\tau_i^{I'})^{-1}(L) \text{ for } i \notin I'. \quad (8)$$

Case 2. $i \in I'$

From definitions of $\Pi_{I'}^I, \tau_i^I$ and $\tau_i^{I'}$ follows

$$\tau_i^I = \tau_i^{I'} \circ \Pi_{I'}^I \text{ for } i \in I'. \quad (9)$$

For $x \in \Pi_{I'}^I((\tau_i^I)^{-1}(L))$ exists $y \in \Sigma_I^*$ with $\tau_i^I(y) \in L$ and $x = \Pi_{I'}^I(y)$. Because of (9) holds

$$\tau_i^{I'}(x) = \tau_i^{I'}(\Pi_{I'}^I(y)) = \tau_i^I(y) \in L,$$

hence, $x \in (\tau_i^{I'})^{-1}(L)$. Therefore,

$$\Pi_{I'}^I((\tau_i^I)^{-1}(L)) \subset (\tau_i^{I'})^{-1}(L) \text{ for } i \in I'. \quad (10)$$

Because of (8), (10) and (7) holds

$$\Pi_{I'}^I((\tau_i^I)^{-1}(L)) = (\tau_i^{I'})^{-1}(L)$$

for $I, I' \in \mathcal{I}$, $I' \subset I$ and $i \in N$. ■

Intersections of system behaviours play an important role concerning uniformity of parameterisation. Therefore, some general properties of intersections of families of sets will be presented.

Let T be a set. A family $f = (f_t)_{t \in T}$ with $f_t \in F$ for each $t \in T$ is formally equivalent to a function $f : T \rightarrow F$ with $f_t := f(t)$.

Let M be a set. A family $f = (f_t)_{t \in T}$ with $f_t \in F = \mathcal{P}(M)$ for each $t \in T$ is called a family of subsets of M .

Let now $T \neq \emptyset$ and f a family of subsets of M . The intersection $\bigcap_{t \in T} f_t$ is defined by

$$\bigcap_{t \in T} f_t = \{m \in M \mid m \in f_t \text{ for each } t \in T\}. \quad (11)$$

If $f = g \circ h$ with $h : T \rightarrow H$ and $g : H \rightarrow F$ then

$$\bigcap_{t \in T} f(t) = \bigcap_{x \in h(T)} g(x). \quad (12)$$

If especially $f = h$ and g is the identity on F , then from (12) follows

$$\bigcap_{t \in T} f(t) = \bigcap_{x \in f(T)} x.$$

For a second family of sets $f' : T' \rightarrow F$ with $f'(T') = f(T)$ follows then

$$\bigcap_{t \in T} f(t) = \bigcap_{t' \in T'} f'(t').$$

In the following we will use family and function notations side by side.

Let $f = (f_t)_{t \in T}$ a family of sets with $f : T \rightarrow F = \mathcal{P}(M)$. If $T = \hat{T} \cup \check{T}$ with $\hat{T} \neq \emptyset$ and $f(\hat{T}) = \{M\}$, then from (11) follows

$$\bigcap_{t \in T} f(t) = \bigcap_{t \in \hat{T}} f(t). \quad (13)$$

Let $\mathcal{E}_{\mathcal{I}} = (E(t, I))_{(t, I) \in T \times \mathcal{I}}$ be an influence structure for \mathcal{I} indexed by T .

For each $I \in \mathcal{I}$ a family of sets

$$\mathcal{E}_{\mathcal{I}}(I) := (E(t, I))_{t \in T}$$

with $E(t, I) = \mathcal{E}_{\mathcal{I}}(I)(t) \in \mathcal{P}(I)$ is defined, and it holds

$$\mathcal{E}_{\mathcal{I}}(I) : T \rightarrow \mathcal{P}(I). \quad (14)$$

From (12) it follows (with $h = \mathcal{E}_{\mathcal{I}}(I)$)

$$\bigcap_{t \in T} (\tau_{E(t, I)}^I)^{-1}(V) = \bigcap_{x \in \mathcal{E}_{\mathcal{I}}(I)(T)} (\tau_x^I)^{-1}(V) \quad (15)$$

for each $V \subset \Sigma^*$ and $I \in \mathcal{I}$.

For each $I \in \mathcal{I}$ holds $\tau_{\emptyset}^I(w) = \varepsilon$ for each $w \in \Sigma_I^*$. It follows,

$$(\tau_{\emptyset}^I)^{-1}(V) = \Sigma_I^* \text{ if } \varepsilon \in V \subset \Sigma^*. \quad (16)$$

Because of (12), (13), (15), and (16)

$$\begin{aligned} \bigcap_{t \in T} (\tau_{E(t, I)}^I)^{-1}(V) &= \bigcap_{x \in \mathcal{E}_{\mathcal{I}}(I)(T_I)} (\tau_x^I)^{-1}(V) \\ &= \bigcap_{t \in T_I} (\tau_{E(t, I)}^I)^{-1}(V) \end{aligned} \quad (17)$$

for each T_I with $\emptyset \neq T_I \subset T$ and $\mathcal{E}_{\mathcal{I}}(I)(T) \setminus \mathcal{E}_{\mathcal{I}}(I)(T_I) \in \{\emptyset, \{\emptyset\}\}$ and $\varepsilon \in V \subset \Sigma^*$.

Each bijection $\iota : I \rightarrow I'$ defines another bijection $\check{\iota} : \mathcal{P}(I) \rightarrow \mathcal{P}(I')$ by

$$\check{\iota}(x) := \{\iota(y) \in I' \mid y \in x\} \text{ for each } x \in \mathcal{P}(I). \quad (18)$$

Lemma 3. Let $\mathcal{E}_{\mathcal{I}} = (E(t, I))_{(t, I) \in T \times \mathcal{I}}$ be an influence structure for \mathcal{I} indexed by T , and let $\mathcal{B}_{\mathcal{I}} = (\mathcal{B}(I, I'))_{(I, I') \in \mathcal{I} \times \mathcal{I}}$ be an isomorphism structure for \mathcal{I} . Let

$$\begin{aligned} &\varepsilon \in V \subset \Sigma^*, \text{ and let } (T_K)_{K \in \mathcal{I}} \text{ be a family} \\ &\text{with } \emptyset \neq T_K \subset T \text{ and} \\ &\mathcal{E}_{\mathcal{I}}(K)(T) \setminus \mathcal{E}_{\mathcal{I}}(K)(T_K) \in \{\emptyset, \{\emptyset\}\} \text{ for each } K \in \mathcal{I}, \\ &\text{such that } \check{\iota}(\mathcal{E}_{\mathcal{I}}(I)(T_I)) = \mathcal{E}_{\mathcal{I}}(I')(T_{I'}) \\ &\text{for each } (I, I') \in \mathcal{I} \times \mathcal{I} \text{ and } \iota \in \mathcal{B}(I, I'), \end{aligned} \quad (19)$$

then

$$\bigcap_{t \in T} (\tau_{E(t, I)}^I)^{-1}(V) = \bigcap_{t \in T_I} (\tau_{E(t, I)}^I)^{-1}(V) \quad (20)$$

for each $I \in \mathcal{I}$, and

$$\iota_{I'}^I \left[\bigcap_{t \in T} (\tau_{E(t, I)}^I)^{-1}(V) \right] = \bigcap_{t \in T} (\tau_{E(t, I')}^{I'})^{-1}(V) \quad (21)$$

for each $(I, I') \in \mathcal{I} \times \mathcal{I}$ and $\iota \in \mathcal{B}(I, I')$.

Proof of (20): Because of (17) from assumption (19) directly follows (20). ■

For the proof of (21) the following property of the homomorphisms τ_K^I is needed:

Let $\iota : I \rightarrow I'$ a bijection and $K \subset I$, then $\tau_{\iota(K)}^{I'} \circ \iota_{I'}^I = \tau_K^I$ and so

$$\tau_{\iota(K)}^{I'} = \tau_K^I \circ (\iota_{I'}^I)^{-1}. \quad (22)$$

Proof of (22):

The elements of $\Sigma_{I'}$ are of the form a_i with $i \in I$ and $a \in \Sigma$. For these elements holds

$$\begin{aligned} \tau_K^I(a_i) &= \begin{cases} a & i \in K \\ \varepsilon & i \in I \setminus K \end{cases} \\ &= \begin{cases} a & \iota(i) \in \iota(K) \\ \varepsilon & \iota(i) \in I' \setminus \iota(K) \end{cases} \\ &= \tau_{\iota(K)}^{I'}(a_{\iota(i)}) = \tau_{\iota(K)}^{I'}(\iota_{I'}^I(a_i)) \end{aligned}$$

which proves (22). ■

Proof of (21): Because of (17) and (22)

$$\begin{aligned} &\iota_{I'}^I \left[\bigcap_{t \in T} (\tau_{E(t, I)}^I)^{-1}(V) \right] \\ &= \iota_{I'}^I \left[\bigcap_{x \in \mathcal{E}_{\mathcal{I}}(I)(T_I)} (\tau_x^I)^{-1}(V) \right] \\ &= ((\iota_{I'}^I)^{-1})^{-1} \left[\bigcap_{x \in \mathcal{E}_{\mathcal{I}}(I)(T_I)} (\tau_x^I)^{-1}(V) \right] \\ &= \bigcap_{x \in \mathcal{E}_{\mathcal{I}}(I)(T_I)} ((\iota_{I'}^I)^{-1})^{-1} [(\tau_x^I)^{-1}(V)] \\ &= \bigcap_{x \in \mathcal{E}_{\mathcal{I}}(I)(T_I)} (\tau_x^I \circ (\iota_{I'}^I)^{-1})^{-1}(V) \\ &= \bigcap_{x \in \mathcal{E}_{\mathcal{I}}(I)(T_I)} (\tau_{\iota(x)}^{I'})^{-1}(V) \\ &= \bigcap_{x \in \mathcal{E}_{\mathcal{I}}(I)(T_I)} (\tau_{\check{\iota}(x)}^{I'})^{-1}(V). \end{aligned} \quad (23)$$

From (12) (with $h = \check{\iota}$) and the assumption (19) follows

$$\begin{aligned} \bigcap_{x \in \mathcal{E}_{\mathcal{I}}(I)(T_I)} (\tau_{\check{\iota}(x)}^{I'})^{-1}(V) &= \bigcap_{x' \in \check{\iota}(\mathcal{E}_{\mathcal{I}}(I)(T_I))} (\tau_{x'}^{I'})^{-1}(V) \\ &= \bigcap_{x' \in \mathcal{E}_{\mathcal{I}}(I')(T_{I'})} (\tau_{x'}^{I'})^{-1}(V). \end{aligned}$$

Furthermore, from (17) follows

$$\bigcap_{x' \in \mathcal{E}_{\mathcal{I}}(I')(T_{I'})} (\tau_{x'}^{I'})^{-1}(V) = \bigcap_{t \in T} (\tau_{E(t, I')}^{I'})^{-1}(V). \quad (24)$$

(23) - (24) prove (21). ■

The case $T = N$, where \mathcal{I} is based on N , allows a simpler sufficient condition for (20) and (21).

Lemma 4. Let \mathcal{I} be a parameter structure based on N , $\mathcal{E}_{\mathcal{I}} = (E(n, I))_{(n, I) \in N \times \mathcal{I}}$ be an influence structure for \mathcal{I} , and let $\mathcal{B}_{\mathcal{I}} = (\mathcal{B}(I, I'))_{(I, I') \in \mathcal{I} \times \mathcal{I}}$ be an isomorphism

structure for \mathcal{I} .

Let $\varepsilon \in V \subset \Sigma^*$, (25a)

for each $I \in \mathcal{I}$ and $n \in N$ let $E(n, I) = \emptyset$,

or it exists an $i_n \in I$ with $E(n, I) = E(i_n, I)$, and (25b)

for each $(I, I') \in \mathcal{I} \times \mathcal{I}$, $\iota \in \mathcal{B}(I, I')$ and $i \in I$ holds

$\iota(E(i, I)) = E(\iota(i), I')$. (25c)

Then

$$\bigcap_{n \in N} (\tau_{E(n, I)}^I)^{-1}(V) = \bigcap_{n \in I} (\tau_{E(n, I)}^I)^{-1}(V) \quad (26)$$

for each $I \in \mathcal{I}$, and

$$\iota_{I'}^I \left[\bigcap_{n \in N} (\tau_{E(n, I)}^I)^{-1}(V) \right] = \bigcap_{n \in N} (\tau_{E(n, I')}^{I'})^{-1}(V) \quad (27)$$

for each $(I, I') \in \mathcal{I} \times \mathcal{I}$ and $\iota \in \mathcal{B}(I, I')$.

Proof: From (25b) follows $\mathcal{E}_{\mathcal{I}}(I)(N) = \mathcal{E}_{\mathcal{I}}(I)(I)$ or $\mathcal{E}_{\mathcal{I}}(I)(N) = \mathcal{E}_{\mathcal{I}}(I)(I) \cup \{\emptyset\}$, so

$$\mathcal{E}_{\mathcal{I}}(I)(N) \setminus \mathcal{E}_{\mathcal{I}}(I)(I) \in \{\emptyset, \{\emptyset\}\} \text{ for each } I \in \mathcal{I}. \quad (28)$$

From (25c) follows

$$\check{\iota}(\mathcal{E}_{\mathcal{I}}(I)(I)) \subset \mathcal{E}_{\mathcal{I}}(I')(I'). \quad (29)$$

Because $\iota : I \rightarrow I'$ is a bijection, for each $i' \in I'$ exists an $i \in I$ with $\iota(i) = i'$. Because of (25c) holds $\check{\iota}(E(i, I)) = E(i', I')$, where $E(i, I) \in \mathcal{E}_{\mathcal{I}}(I)(I)$. From this follows

$$\mathcal{E}_{\mathcal{I}}(I')(I') \subset \check{\iota}(\mathcal{E}_{\mathcal{I}}(I)(I)). \quad (30)$$

Because of (28) - (30), with $T = N$ and $(T_I)_{I \in \mathcal{I}} = (I)_{I \in \mathcal{I}}$,

$$(25a) - (25c) \text{ implies } (19).$$

Example 2 (Example 1 (continued)). Let \mathcal{I} be a parameter structure based on N and $\mathcal{B}_{\mathcal{I}} = (\mathcal{B}(I, I'))_{(I, I') \in \mathcal{I} \times \mathcal{I}}$ be an isomorphism structure for \mathcal{I} . Then $\dot{\mathcal{E}}_{\mathcal{I}}$ satisfies (25b) and (25c).

So for $\varepsilon \in L \subset \Sigma^*$ Lemma 4 implies

$$\begin{aligned} \bigcap_{n \in N} (\tau_n^I)^{-1}(L) &= \bigcap_{n \in I} (\tau_n^I)^{-1}(L) \text{ for each } I \in \mathcal{I} \text{ and} \\ \iota_{I'}^I \left[\bigcap_{n \in N} (\tau_n^I)^{-1}(L) \right] &= \bigcap_{n \in N} (\tau_n^{I'})^{-1}(L) \end{aligned} \quad (31)$$

for each $(I, I') \in \mathcal{I} \times \mathcal{I}$ and $\iota \in \mathcal{B}(I, I')$.

Now Lemma 2 together with (31) proves Theorem 2.

(32)

Because of $\tau_n^I = \tau_{\dot{E}(n, I)}^I$ for $I \in \mathcal{I}$ and $n \in N$, (31) and the definitions of $(\dot{\mathcal{L}}(L)_I)_{I \in \mathcal{I}}$ and $(\mathcal{L}(L, \mathcal{E}_{\mathcal{I}}, V)_I)_{I \in \mathcal{I}}$ imply

$$\begin{aligned} \dot{\mathcal{L}}(L)_I &= \bigcap_{n \in I} (\tau_n^I)^{-1}(L) = \bigcap_{n \in I} (\tau_n^I)^{-1}(L) \cap \bigcap_{n \in I} (\tau_n^I)^{-1}(V) \\ &= \dot{\mathcal{L}}(L)_I \cap \bigcap_{n \in N} (\tau_n^I)^{-1}(V) \\ &= \dot{\mathcal{L}}(L)_I \cap \bigcap_{n \in N} (\tau_{\dot{E}(n, I)}^I)^{-1}(V) \\ &= \mathcal{L}(L, \dot{\mathcal{E}}_I, V)_I \end{aligned} \quad (33)$$

for $I \in \mathcal{I}$ and $V \supset L$.

(33) gives a representation of $(\dot{\mathcal{L}}(L)_I)_{I \in \mathcal{I}}$ in terms of $(\mathcal{L}(L, \mathcal{E}_{\mathcal{I}}, V)_I)_{I \in \mathcal{I}}$.

For the following theorems please remember that by the general definition of $\mathcal{L}(L, \mathcal{E}_{\mathcal{I}}, V)_I$ it is assumed that $\emptyset \neq L \subset V$ and L, V are prefix closed. This implies $\varepsilon \in L \subset V$.

Lemma 5. Let \mathcal{I} be a parameter structure, $\mathcal{E}_{\mathcal{I}}$ an influence structure for \mathcal{I} indexed by T and $\mathcal{B}_{\mathcal{I}}$ an isomorphism structure for \mathcal{I} .

Assuming (2) and (19), then

$$(\mathcal{L}(L, \mathcal{E}_{\mathcal{I}}, V)_I)_{I \in \mathcal{I}}$$

is a scalable systems with respect to $\mathcal{B}_{\mathcal{I}}$.

$$\text{It holds } \mathcal{L}(L, \mathcal{E}_{\mathcal{I}}, V)_I = \dot{\mathcal{L}}(L)_I \cap \bigcap_{n \in T_I} (\tau_{E(n, I)}^I)^{-1}(V)$$

for each $I \in \mathcal{I}$.

Proof: By Theorem 2, $(\dot{\mathcal{L}}(L)_I)_{I \in \mathcal{I}}$ is a scalable system with respect to $\mathcal{B}_{\mathcal{I}}$. By Lemma 1 and 3 (21)

$$\left(\bigcap_{t \in T} (\tau_{E(t, I)}^I)^{-1}(V) \right)_{I \in \mathcal{I}}$$

is a scalable system with respect to $\mathcal{B}_{\mathcal{I}}$ too. Now part (ii) of the intersection theorem proves $(\mathcal{L}(L, \mathcal{E}_{\mathcal{I}}, V)_I)_{I \in \mathcal{I}}$ to be a scalable system with respect to $\mathcal{B}_{\mathcal{I}}$. Lemma 3 (20) completes the proof of Lemma 5. ■

Using Lemma 4 instead of Lemma 3 proves the following.

Theorem 3 (construction condition for scalable systems). By the assumptions of Lemma 4 and (2) with $T = N$, $(\mathcal{L}(L, \mathcal{E}_{\mathcal{I}}, V)_I)_{I \in \mathcal{I}}$ is a scalable system with respect to $\mathcal{B}_{\mathcal{I}}$. It holds $\mathcal{L}(L, \mathcal{E}_{\mathcal{I}}, V)_I = \dot{\mathcal{L}}(L)_I \cap \bigcap_{n \in I} (\tau_{E(n, I)}^I)^{-1}(V)$.

Remark 1. It can be shown that in $\text{SP}(L, V) \mathbb{N}$ can be replaced by each countable infinite set. Let therefore N' be another set and $\iota : \mathbb{N} \rightarrow N'$ a bijection. $\iota_{N'}^{\mathbb{N}} : \Sigma_{\mathbb{N}}^* \rightarrow \Sigma_{N'}^*$ is the isomorphism defined as in the definition of isomorphism structure. It now holds

$$\Theta^{\mathbb{N}} = \Theta^{N'} \circ \iota_{N'}^{\mathbb{N}} \text{ and } \tau_n^{\mathbb{N}} = \tau_{\iota(n)}^{N'} \circ \iota_{N'}^{\mathbb{N}} \quad (34)$$

for each $n \in \mathbb{N}$. Furthermore,

$$\iota_{N'}^{\mathbb{N}} \circ \Pi_K^{\mathbb{N}} = \Pi_{\iota(K)}^{N'} \circ \iota_{N'}^{\mathbb{N}} \quad (35)$$

for each $K \subset \mathbb{N}$. From (34) and commutativity of intersection now

$$\begin{aligned} & \left(\bigcap_{n \in \mathbb{N}} (\tau_n^{\mathbb{N}})^{-1}(L) \right) \cap (\Theta^{\mathbb{N}})^{-1}(V) = \\ & = (\iota_{N'}^{\mathbb{N}})^{-1} \left[\left(\bigcap_{n \in \mathbb{N}} (\tau_{\iota(n)}^{N'})^{-1}(L) \right) \cap (\Theta^{N'})^{-1}(V) \right] \\ & = (\iota_{N'}^{\mathbb{N}})^{-1} \left[\left(\bigcap_{n' \in N'} (\tau_{n'}^{N'})^{-1}(L) \right) \cap (\Theta^{N'})^{-1}(V) \right]. \end{aligned} \quad (36)$$

By (35),

$$\Pi_K^{\mathbb{N}} \circ (\iota_{N'}^{\mathbb{N}})^{-1} = (\iota_{N'}^{\mathbb{N}})^{-1} \circ \Pi_{\iota(K)}^{N'}. \quad (37)$$

Because of (36) and (37)

$$\begin{aligned} & \Pi_K^{\mathbb{N}} \left[\left(\bigcap_{n \in \mathbb{N}} (\tau_n^{\mathbb{N}})^{-1}(L) \right) \cap (\Theta^{\mathbb{N}})^{-1}(V) \right] = \\ & = (\iota_{N'}^{\mathbb{N}})^{-1} \left(\Pi_{\iota(K)}^{N'} \left[\left(\bigcap_{n' \in N'} (\tau_{n'}^{N'})^{-1}(L) \right) \cap (\Theta^{N'})^{-1}(V) \right] \right). \end{aligned} \quad (38)$$

From

$$\Pi_K^{\mathbb{N}} \left[\left(\bigcap_{n \in \mathbb{N}} (\tau_n^{\mathbb{N}})^{-1}(L) \right) \cap (\Theta^{\mathbb{N}})^{-1}(V) \right] \subset (\Theta^{\mathbb{N}})^{-1}(V)$$

now follows

$$\begin{aligned} & \Pi_{\iota(K)}^{N'} \left[\left(\bigcap_{n' \in N'} (\tau_{n'}^{N'})^{-1}(L) \right) \cap (\Theta^{N'})^{-1}(V) \right] \\ & \subset \iota_{N'}^{\mathbb{N}} \left((\Theta^{\mathbb{N}})^{-1}(V) \right). \end{aligned} \quad (39)$$

Because of (34) $\Theta^{\mathbb{N}} \circ (\iota_{N'}^{\mathbb{N}})^{-1} = \Theta^{N'}$ and so

$$(\Theta^{N'})^{-1}(V) = \iota_{N'}^{\mathbb{N}} \left((\Theta^{\mathbb{N}})^{-1}(V) \right).$$

Therefore, from (39) follows

$$\Pi_{\iota(K)}^{N'} \left[\left(\bigcap_{n' \in N'} (\tau_{n'}^{N'})^{-1}(L) \right) \cap (\Theta^{N'})^{-1}(V) \right] \subset (\Theta^{N'})^{-1}(V). \quad (40)$$

Because for each $\emptyset \neq K' \subset N'$ it exists an $\emptyset \neq K \subset \mathbb{N}$ with $K' = \iota(K)$, by $\text{SP}(L, V)$, we get for each $\emptyset \neq K \subset \mathbb{N}$ a corresponding inclusion with N' replacing \mathbb{N} and K' for K .

Lemma 6. *The assumptions of Lemma 1 and Lemma 2 together with $\text{SP}(L, V)$ imply that $(X(L, V, t))_{I \in \mathcal{I}}$ with*

$$X(L, V, t)_I := \bigcap_{n \in N} (\tau_n^I)^{-1}(L) \cap (\tau_{E(t, I)}^I)^{-1}(V)$$

is a self-similar monotonic parameterised system for each $t \in T$.

Proof: By Lemma 1 and Lemma 2, $((\tau_{E(t, I)}^I)^{-1}(V))_{I \in \mathcal{I}}$ and $(\bigcap_{n \in N} (\tau_n^I)^{-1}(L))_{I \in \mathcal{I}}$ are monotonic parameterised systems. So by the intersection theorem $(X(L, V, t))_{I \in \mathcal{I}}$ is a monotonic parameterised system for each $t \in T$. Therefore

$$X(L, V, t)_{I'} = \Pi_{I'}^I(X(L, V, t)_{I'}) \subset \Pi_{I'}^I(X(L, V, t)_I)$$

for each $I, I' \in \mathcal{I}$ with $I' \subset I$. So the proof of self-similarity can be reduced to the proof of

$$\Pi_{I'}^I(X(L, V, t)_I) \subset X(L, V, t)_{I'} \quad (41)$$

for each $t \in T$ and $I, I' \in \mathcal{I}$ with $I' \subset I$.

Because by Lemma 2

$$\left(\bigcap_{n \in N} (\tau_n^I)^{-1}(L) \right)_{I \in \mathcal{I}}$$

is self-similar, it holds

$$\Pi_{I'}^I(X(L, V, t)_I) \subset \Pi_{I'}^I \left(\bigcap_{n \in N} (\tau_n^I)^{-1}(L) \right) = \bigcap_{n \in N} (\tau_n^I)^{-1}(L).$$

So the proof of (41) can be reduced to the proof of

$$\Pi_{I'}^I \left[\left(\bigcap_{n \in N} (\tau_n^I)^{-1}(L) \right) \cap (\tau_{E(t, I)}^I)^{-1}(V) \right] \subset (\tau_{E(t, I')}^{I'})^{-1}(V) \quad (42)$$

for each $t \in T$ and $I, I' \in \mathcal{I}$ with $I' \subset I$.

For each $w \in \left(\bigcap_{n \in N} (\tau_n^I)^{-1}(L) \right) \cap (\tau_{E(t, I)}^I)^{-1}(V)$ exists a $r \in \mathbb{N}$ and $u_i \in \Sigma_{E(t, I)}^{*n}$ for $1 \leq i \leq r$ and $v_i \in \Sigma_{I \setminus E(t, I)}^*$ for $1 \leq i \leq r$ with $w = u_1 v_1 u_2 v_2 \dots u_r v_r$.

Note that $\Sigma_\emptyset := \emptyset$ and $\emptyset^* = \{\varepsilon\}$.

Because $u_1 u_2 \dots u_r \in \Sigma_{E(t, I)}^*$ and $v_1 v_2 \dots v_r \in \Sigma_{I \setminus E(t, I)}^*$ holds

$$\begin{aligned} \Theta^N(u_1 u_2 \dots u_r) &= \tau_{E(t, I)}^I(u_1 u_2 \dots u_r) \\ &= \tau_{E(t, I)}^I(w) \in V. \end{aligned} \quad (43)$$

With the same argumentation holds

$$\tau_n^N(u_1 u_2 \dots u_r) = \tau_n^I(u_1 u_2 \dots u_r) = \tau_n^I(w) \in L \quad (44)$$

for $n \in E(t, I)$ and

$$\tau_n^N(u_1 u_2 \dots u_r) = \varepsilon \in L \quad (45)$$

for $n \in N \setminus E(t, I)$. With (43) - (45) now

$$u_1 u_2 \dots u_r \in \left(\bigcap_{n \in N} (\tau_n^N)^{-1}(L) \right) \cap (\Theta^N)^{-1}(V),$$

and on behalf of precondition $\text{SP}(L, V)$ holds

$$\begin{aligned} \Pi_{I'}^N(u_1 u_2 \dots u_r) &= \Pi_{I' \cap E(t, I)}^{E(t, I)}(u_1 u_2 \dots u_r) \\ &\in \Sigma_{I' \cap E(t, I)}^* \cap (\Theta^N)^{-1}(V). \end{aligned} \quad (46)$$

Furthermore,

$$\begin{aligned} \Pi_{I'}^I(w) &= \Pi_{I'}^I(u_1 v_1 u_2 v_2 \dots u_r v_r) \\ &= \Pi_{I' \cap E(t, I)}^{E(t, I)}(u_1) \Pi_{I' \setminus E(t, I)}^{I \setminus E(t, I)}(v_1) \dots \\ &\quad \Pi_{I' \cap E(t, I)}^{E(t, I)}(u_r) \Pi_{I' \setminus E(t, I)}^{I \setminus E(t, I)}(v_r). \end{aligned} \quad (47)$$

Because of (2), $E(t, I') \subset E(t, I)$ and so $I' \setminus E(t, I) \subset I' \setminus E(t, I')$ and thus

$$\tau_{E(t, I')}^{I'}(\Pi_{I' \setminus E(t, I)}^{I \setminus E(t, I)}(v_i)) = \varepsilon$$

for $1 \leq i \leq r$. With (2) and (47) it follows

$$\tau_{E(t, I')}^{I'}(\Pi_{I'}^I(w)) = \tau_{E(t, I')}^{I'}(\Pi_{E(t, I')}^{E(t, I)}(u_1 \dots u_r)). \quad (48)$$

Because $\tau_{E(t,I')}^{I'}(x) = \Theta^N(x)$ for each $x \in \Sigma_{E(t,I')}^*$ now on behalf of (48), (2), and (46)

$$\tau_{E(t,I')}^{I'}(\Pi_{I'}^I(w)) = \Theta^N(\Pi_{E(t,I')}^{E(t,I)}(u_1 \dots u_r)) \in V,$$

and thus $\Pi_{I'}^I(w) \in (\tau_{E(t,I')}^{I'})^{-1}(V)$. This proves (42) and completes the proof of Lemma 6. ■

Because of the idempotence of intersection

$$\begin{aligned} & \bigcap_{n \in N} (\tau_n^I)^{-1}(L) \cap \bigcap_{t \in T} (\tau_{E(t,I)}^I)^{-1}(V) \\ &= \bigcap_{t \in T} \left[\bigcap_{n \in N} (\tau_n^I)^{-1}(L) \cap (\tau_{E(t,I)}^I)^{-1}(V) \right]. \end{aligned}$$

Now the intersection theorem and Lemma 6 imply

Lemma 7. *If $\text{SP}(L, V)$, then by the assumptions of Lemma 1 and 2*

$$\left[\bigcap_{n \in N} (\tau_n^I)^{-1}(L) \cap \bigcap_{t \in T} (\tau_{E(t,I)}^I)^{-1}(V) \right]_{I \in \mathcal{I}}$$

is a self-similar monotonic parameterised system.

Combining Lemma 7 with Lemma 5 or Theorem 3 imply

Theorem 4 (construction condition for well-behaved scalable systems). *By the assumptions of Lemma 5 or Theorem 3 together with $\text{SP}(L, V)$*

$$(\mathcal{L}(L, \mathcal{E}_{\mathcal{I}}, V)_I)_{I \in \mathcal{I}}$$

is a well-behaved scalable system.

Theorem 5 (inverse abstraction theorem). *Let $\varphi : \Sigma^* \rightarrow \Phi^*$ be an alphabetic homomorphism and $W, X \subset \Phi^*$, then*

$$\text{SP}(W, X) \text{ implies } \text{SP}(\varphi^{-1}(W), \varphi^{-1}(X)).$$

Proof of Theorem 5:

Let K be a non-empty set. Each alphabetic homomorphism $\varphi : \Sigma^* \rightarrow \Phi^*$ defines a homomorphism $\varphi^K : \Sigma_K^* \rightarrow \Phi_K^*$ by

$$\varphi^K(a_n) := (\varphi(a))_n \text{ for } a_n \in \Sigma_K, \text{ where } (\varepsilon)_n = \varepsilon. \quad (49)$$

If $\bar{\tau}_n^K : \Phi_K^* \rightarrow \Phi$ and $\bar{\Theta}^K : \Phi_K^* \rightarrow \Phi$ are defined analogous to $\bar{\tau}_n^K$ and $\bar{\Theta}^K$, then

$$\varphi \circ \bar{\tau}_n^K = \bar{\tau}_n^K \circ \varphi^K, \text{ and } \varphi \circ \bar{\Theta}^K = \bar{\Theta}^K \circ \varphi^K. \quad (50)$$

Let now N be an infinite countable set. Because of (50), for $W, X \subset \Phi^*$

$$\begin{aligned} & \left(\bigcap_{n \in N} (\tau_n^N)^{-1}(\varphi^{-1}(W)) \right) \cap (\bar{\Theta}^N)^{-1}(\varphi^{-1}(X)) \\ &= (\varphi^N)^{-1} \left[\left(\bigcap_{n \in N} (\bar{\tau}_n^N)^{-1}(W) \right) \cap (\bar{\Theta}^N)^{-1}(X) \right]. \end{aligned} \quad (51)$$

Because of $\varphi^K(w) = \varphi^N(w)$ for $w \in \Sigma_K^* \subset \Sigma_N^*$ and $\emptyset \neq K \subset N$

$$(\varphi^K)^{-1}(Z) \subset (\varphi^N)^{-1}(Z) \text{ for } Z \subset \Phi_K^*. \quad (52)$$

If now $\text{SP}(W, X)$, and

$$\Pi_K^N[(\varphi^N)^{-1}(Y)] = (\varphi^K)^{-1}(\bar{\Pi}_K^N[Y]) \quad (53)$$

for $Y \subset \Phi_N^*$ and $\emptyset \neq K \subset N$, where $\bar{\Pi}_K^N : \Phi_N^* \rightarrow \Phi_K^*$ is defined analogous to Π_K^N , then follows (with (50) - (53))

$$\begin{aligned} & \Pi_K^N \left[\left(\bigcap_{n \in N} (\tau_n^N)^{-1}(\varphi^{-1}(W)) \right) \cap (\bar{\Theta}^N)^{-1}(\varphi^{-1}(X)) \right] \\ &= (\varphi^K)^{-1} \left[\bar{\Pi}_K^N \left[\left(\bigcap_{n \in N} (\bar{\tau}_n^N)^{-1}(W) \right) \cap (\bar{\Theta}^N)^{-1}(X) \right] \right] \\ &\subset (\varphi^K)^{-1} \left[(\bar{\Theta}^N)^{-1}(X) \right] \subset (\varphi^N)^{-1} \left[(\bar{\Theta}^N)^{-1}(X) \right] \\ &= (\bar{\Theta}^N)^{-1}(\varphi^{-1}(X)). \end{aligned} \quad (54)$$

With (54)

$$\text{SP}(\varphi^{-1}(W), \varphi^{-1}(X)) \text{ follows from } \text{SP}(W, X), \quad (55)$$

if (53) holds.

It remains to show (53). For the proof of (53) it is sufficient to prove

$$\Pi_K^N((\varphi^N)^{-1}(y)) = (\varphi^K)^{-1}(\bar{\Pi}_K^N(y)) \quad (56)$$

for each $y \in \Phi_N^*$, because of

$$\Pi_K^N((\varphi^N)^{-1}(Y)) = \bigcup_{y \in Y} \Pi_K^N((\varphi^N)^{-1}(y))$$

and

$$(\varphi^K)^{-1}(\bar{\Pi}_K^N(Y)) = \bigcup_{y \in Y} (\varphi^K)^{-1}(\bar{\Pi}_K^N(y)).$$

Here, for $f : A \rightarrow B$ and $b \in B$ we use the convention

$$f^{-1}(b) = f^{-1}(\{b\}).$$

With $Y = \{y\}$ (56) is also necessary for (53), and so it is equivalent to (53).

Definition 1 ((general) projection). *For arbitrary alphabets Δ and Δ' with $\Delta' \subset \Delta$ general projections $\pi_{\Delta'}^{\Delta} : \Delta^* \rightarrow \Delta'^*$ are defined by*

$$\pi_{\Delta'}^{\Delta}(a) := \begin{cases} a & | \quad a \in \Delta' \\ \varepsilon & | \quad a \in \Delta \setminus \Delta' \end{cases}. \quad (57)$$

In this terminology the projections

$$\Pi_K^N : \Sigma_N^* \rightarrow \Sigma_K^* \text{ and } \bar{\Pi}_K^N : \Phi_N^* \rightarrow \Phi_K^*$$

considered until now are special cases, which we call *parameter-projections*. It holds

$$\Pi_K^N = \pi_{\Sigma_K^*}^{\Sigma_N^*} \text{ and } \bar{\Pi}_K^N = \pi_{\Phi_K^*}^{\Phi_N^*}. \quad (58)$$

Because of the different notations, in general we just use the term *projection* for both cases.

We now consider the equation (56) for the special case, where $\varphi : \Sigma^* \rightarrow \Phi^*$ is a projection, that is $\varphi = \pi_{\Phi^*}^{\Sigma^*}$ with $\Phi \subset \Sigma$. In this case also $\varphi^N : \Sigma_N^* \rightarrow \Phi_N^*$ is a projection, with

$$\varphi^N = \pi_{\Phi_N^*}^{\Sigma_N^*}. \quad (59)$$

Lemma 8 (projection-lemma).

Let Δ be an alphabet, $\Delta' \subset \Delta$, $\Gamma \subset \Delta$ and $\Gamma' = \Delta' \cap \Gamma$, then

$$\pi_{\Delta'}^{\Delta}((\pi_{\Gamma}^{\Delta})^{-1}(y)) = (\pi_{\Gamma'}^{\Delta'})^{-1}(\pi_{\Delta'}^{\Delta}(y))$$

for each $y \in \Gamma^*$.

Proof: Let $y \in \Gamma^*$. We show

$$\pi_{\Gamma'}^{\Delta'}(\pi_{\Delta'}^{\Delta}(z)) = \pi_{\Delta'}^{\Delta}(y) \text{ for each } z \in (\pi_{\Gamma}^{\Delta})^{-1}(y) \quad (60)$$

and we show that

$$\begin{aligned} &\text{for each } u \in (\pi_{\Gamma'}^{\Delta'})^{-1}(\pi_{\Delta'}^{\Delta}(y)) \text{ there exists a} \\ &v \in (\pi_{\Gamma}^{\Delta})^{-1}(y) \text{ such that } \pi_{\Delta'}^{\Delta}(v) = u. \end{aligned} \quad (61)$$

From (60) it follows that

$$\pi_{\Delta'}^{\Delta}((\pi_{\Gamma}^{\Delta})^{-1}(y)) \subset (\pi_{\Gamma'}^{\Delta'})^{-1}(\pi_{\Delta'}^{\Delta}(y))$$

and from (61) it follows that

$$(\pi_{\Gamma'}^{\Delta'})^{-1}(\pi_{\Delta'}^{\Delta}(y)) \subset \pi_{\Delta'}^{\Delta}((\pi_{\Gamma}^{\Delta})^{-1}(y)),$$

which in turn proves Lemma 8.

Proof of (60): By definition of π_{Γ}^{Δ} , $\pi_{\Gamma'}^{\Delta'}$ and $\pi_{\Delta'}^{\Delta}$, follows

$$\pi_{\Gamma'}^{\Delta'}(\pi_{\Delta'}^{\Delta}(z)) = \pi_{\Delta'}^{\Delta}(\pi_{\Gamma}^{\Delta}(z))$$

for each $z \in \Delta^*$ and therewith (60).

Proof of (61) by induction on $y \in \Gamma^*$:

Induction base. Let $y = \varepsilon$, then $u \in (\Delta' \setminus \Gamma')^*$ for each $u \in (\pi_{\Gamma'}^{\Delta'})^{-1}(\pi_{\Delta'}^{\Delta}(y))$. From this follows

$$\pi_{\Delta'}^{\Delta}(v) = u \text{ with } v := u \in (\pi_{\Gamma}^{\Delta})^{-1}(\varepsilon).$$

Induction step. Let $y = \hat{y}\hat{y}$ with $\hat{y} \in \Gamma^*$ and $\hat{y} \in \Gamma$.

Case 1: $\hat{y} \in \Gamma \setminus \Gamma' = \Gamma \cap (\Delta \setminus \Delta')$

Then

$$(\pi_{\Gamma'}^{\Delta'})^{-1}(\pi_{\Delta'}^{\Delta}(y)) = (\pi_{\Gamma'}^{\Delta'})^{-1}(\pi_{\Delta'}^{\Delta}(\hat{y})).$$

By induction hypothesis then for each $u \in (\pi_{\Gamma'}^{\Delta'})^{-1}(\pi_{\Delta'}^{\Delta}(y))$ it exists $\hat{v} \in (\pi_{\Gamma}^{\Delta})^{-1}(\hat{y})$ such that $\pi_{\Delta'}^{\Delta}(\hat{v}) = u$.

With $v := \hat{v}\hat{y}$ holds $\pi_{\Delta'}^{\Delta}(\hat{v}\hat{y}) = \hat{y}\hat{y} = y$ and hence

$$v \in (\pi_{\Gamma}^{\Delta})^{-1}(y) \text{ and } \pi_{\Delta'}^{\Delta}(v) = \pi_{\Delta'}^{\Delta}(\hat{v})\hat{y} = u.$$

Case 2: $\hat{y} \in \Gamma' \subset \Delta'$

Then $\pi_{\Delta'}^{\Delta}(y) = \pi_{\Delta'}^{\Delta}(\hat{y})\hat{y}$. Therefore, each $u \in (\pi_{\Gamma'}^{\Delta'})^{-1}(\pi_{\Delta'}^{\Delta}(y))$ can be departed into $u = \hat{u}\hat{y}\hat{u}$ with $\hat{u} \in (\pi_{\Gamma'}^{\Delta'})^{-1}(\pi_{\Delta'}^{\Delta}(\hat{y}))$ and $\hat{u} \in (\Delta' \setminus \Gamma')^*$.

By induction hypothesis then exists $\hat{v} \in (\pi_{\Gamma}^{\Delta})^{-1}(\hat{y})$ such that $\pi_{\Delta'}^{\Delta}(\hat{v}) = \hat{y}$.

With $v := \hat{v}\hat{y}\hat{u}$ holds $\pi_{\Delta'}^{\Delta}(\hat{v}\hat{y}\hat{u}) = \hat{y}\hat{y} = y$ and hence

$$v \in (\pi_{\Gamma}^{\Delta})^{-1}(y) \text{ and } \pi_{\Delta'}^{\Delta}(v) = \pi_{\Delta'}^{\Delta}(\hat{v})\hat{y}\hat{u} = \hat{y}\hat{y}\hat{u} = u.$$

This completes the proof of (61). ■

For $y \in \Gamma^*$ holds

$$\pi_{\Delta'}^{\Delta}(y) = \pi_{\Delta' \cap \Gamma}^{\Gamma}(y) = \pi_{\Gamma'}^{\Gamma}(y).$$

Therewith, from Lemma 8 follows

$$\pi_{\Delta'}^{\Delta}((\pi_{\Gamma}^{\Delta})^{-1}(y)) = (\pi_{\Gamma'}^{\Delta'})^{-1}(\pi_{\Gamma'}^{\Gamma}(y)) \text{ for each } y \in \Gamma^*. \quad (62)$$

For $\emptyset \neq K \subset N$, $\Phi \subset \Sigma$, $\Delta := \Sigma_N$, $\Delta' := \Sigma_K$, and $\Gamma := \Phi_N$ holds $\Gamma' = \Delta' \cap \Gamma = \Phi_K$.

Assuming $\varphi = \pi_{\Phi}^{\Sigma}$, which implies $\varphi^K = \pi_{\Phi_K}^{\Sigma_K}$, then from (62) (with (58) and (59)), follows

$$\Pi_K^N((\varphi^N)^{-1}(y)) = (\varphi^K)^{-1}(\bar{\Pi}_K^N(y))$$

for $y \in \Phi_N^*$, and so (56). With this,

premise (53) is fulfilled for (55), when φ is a projection, (63)

which proves Theorem 5 for projections.

Definition 2 (strict alphabetic homomorphism). *Let Σ , Φ alphabets, and $\varphi : \Sigma^* \rightarrow \Phi^*$ a homomorphism. Then φ is called alphabetic, if $\varphi(\Sigma) \subset \Phi \cup \{\varepsilon\}$, and φ is called strict alphabetic, if $\varphi(\Sigma) \subset \Phi$.*

Each alphabetic homomorphism $\varphi : \Sigma^* \rightarrow \Phi^*$ is the composition of a projection with a strict alphabetic homomorphism, more precisely,

$$\varphi = \varphi_S \circ \pi_{\varphi^{-1}(\Phi) \cap \Sigma}^{\Sigma}, \quad (64)$$

where $\varphi_S : (\varphi^{-1}(\Phi) \cap \Sigma)^* \rightarrow \Phi^*$ is the strict alphabetic homomorphism defined by

$$\varphi_S(a) := \varphi(a) \text{ for } a \in \varphi^{-1}(\Phi) \cap \Sigma.$$

For $W, X \subset \Phi^*$ and $\varphi : \Sigma^* \rightarrow \Phi^*$ alphabetic (64) implies

$$\begin{aligned} \varphi^{-1}(W) &= (\pi_{\varphi^{-1}(\Phi) \cap \Sigma}^{\Sigma})^{-1}((\varphi_S)^{-1}(W)) \text{ and} \\ \varphi^{-1}(X) &= (\pi_{\varphi^{-1}(\Phi) \cap \Sigma}^{\Sigma})^{-1}((\varphi_S)^{-1}(X)). \end{aligned} \quad (65)$$

Now with (63) and (65) it remains to prove Theorem 5 for strict alphabetic homomorphisms. This will be done by Lemma 9, which proves (56) for strict alphabetic homomorphisms.

Lemma 9. *Let $\varphi : \Sigma^* \rightarrow \Phi^*$ be a strict alphabetic homomorphism, then for all $y \in \Phi_N^*$ and $\emptyset \neq K \subset N$ holds*

$$\Pi_K^N((\varphi^N)^{-1}(y)) = (\varphi^K)^{-1}(\bar{\Pi}_K^N(y)).$$

Proof: Proof by induction on y .

Induction basis: $y = \varepsilon$

Because φ^N is strict alphabetic

$$(\varphi^N)^{-1}(\varepsilon) = \{\varepsilon\} \text{ and so } \Pi_K^N((\varphi^N)^{-1}(\varepsilon)) = \{\varepsilon\}.$$

For the same reason

$$(\varphi^K)^{-1}(\bar{\Pi}_K^N(\varepsilon)) = (\varphi^K)^{-1}(\varepsilon) = \{\varepsilon\}.$$

Induction step: Let $y = y'a_t$ with $a_t \in \Phi_N$, where $a \in \Phi$ and $t \in N$. Because φ^N is alphabetic, it holds

$$(\varphi^N)^{-1}(y'a_t) = ((\varphi^N)^{-1}(y'))((\varphi^N)^{-1}(a_t)),$$

and so

$$\Pi_K^N((\varphi^N)^{-1}(y'a_t)) = \Pi_K^N((\varphi^N)^{-1}(y'))\Pi_K^N((\varphi^N)^{-1}(a_t)).$$

Also holds

$$(\varphi^K)^{-1}(\bar{\Pi}_K^N(y'a_t)) = (\varphi^K)^{-1}(\bar{\Pi}_K^N(y'))(\varphi^K)^{-1}(\bar{\Pi}_K^N(a_t)).$$

According to the induction hypothesis, it holds

$$\Pi_K^N((\varphi^N)^{-1}(y')) = (\varphi^K)^{-1}(\bar{\Pi}_K^N(y')).$$

Therefore, it remains to show

$$\Pi_K^N((\varphi^N)^{-1}(a_t)) = (\varphi^K)^{-1}(\bar{\Pi}_K^N(a_t)).$$

Case 1: $t \notin K$

Because φ^N is strict alphabetic, it holds $(\varphi^N)^{-1}(a_t) \subset \Sigma_t$,
so

$$\Pi_K^N((\varphi^N)^{-1}(a_t)) = \{\varepsilon\}.$$

Additionally holds $\bar{\Pi}_K^N(a_t) = \varepsilon$, and therewith

$$(\varphi^K)^{-1}(\bar{\Pi}_K^N(a_t)) = \{\varepsilon\},$$

because φ^K is strict alphabetic.

Case 2: $t \in K$

Because φ^N is strict alphabetic, it holds

$$(\varphi^N)^{-1}(a_t) = \{b_t \in \Sigma_t \mid \varphi(b) = a\},$$

and therewith

$$\Pi_K^N((\varphi^N)^{-1}(a_t)) = \{b_t \in \Sigma_t \mid \varphi(b) = a\}.$$

$\bar{\Pi}_K^N(a_t) = a_t$ and therewith

$$(\varphi^K)^{-1}(\bar{\Pi}_K^N(a_t)) = \{b_t \in \Sigma_t \mid \varphi(b) = a\},$$

because φ^K is strict alphabetic. This completes the proof
of Lemma 9. ■

This completes the proof of Theorem 5. ■